Labelled Modes A New Notation for Tensor Networks

Simon Etter

University of Warwick

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Introduction

Mathematician's style goals

- Precise.
- Concise.
- General.

The tensor network notation dilemma

- ► Tensor trains: precise & concise, but not general.
- Hierarchical Tucker: general, but not precise or concise.

Introduction

Furthermore

- Even hierarchical Tucker not really general: Requires rooted binary tree, free modes only in leaves.
- Even tensor trains not really precise and concise:

$$x(i_1,\ldots,i_d)=\sum_{\alpha_1}\ldots\sum_{\alpha_{d-1}}x_1(i_1,\alpha_1)x_2(\alpha_1,i_2,\alpha_2)\ldots x_d(\alpha_{d-1},i_d).$$

```
> Coding is hard:
    x = x[1]
    n = n[1]
    for k = 2:d
        x = reshape(x, (n,r[k])) *
            reshape(x[k], (r[k],n[k]*r[k+1]))
        n *= n[k]
end
```

Introduction

Example

- Scientist counting fraction of species with particular features.
- ▶ Natural representation: p(#extremities = x, #eyes = y,...)
- Tensor representation: $p(i_1, \ldots, i_d)$

Example

- Consider two tensors $A \in \mathbb{R}^{4 \times 5 \times 6}$, $B \in \mathbb{R}^{5 \times 7 \times 8}$.
- Only way to contract: 2nd mode of A with 1st mode of B.
- Order of modes in result follows from mode sizes.
- Hence, C = AB contains same information as

$$C(a,c,d,e) = \sum_{b} A(a,b,c) B(b,d,e).$$

Labelled Modes Tensors

Generalised tuple

Let D be a finite set and $(A_k)_{k\in D}$, a family of sets.

• A tuple
$$t \in \bigotimes_{k \in D} A_k$$
 is a function

$$t:D
ightarrowigcup_{k\in D}A_k$$
 such that $t_k\in A_k.$

• Traditional tuples:
$$D = \{1, \ldots, n\}$$
.

Tensor

A *tensor* x is a function mapping integer tuples to scalars,

$$\mathbb{K}(D) := \left\{ x : \bigotimes_{k \in D} [n_k] \to \mathbb{K} \right\}, \text{ where } [n_k] := \{1, \ldots, n_k\}.$$

This construction is natural and intuitive!

Recall example of scientist counting species.

- ▶ Data naturally represented as $p \in \mathbb{R}(\{\# \text{extremities}, \# \text{eyes}\})$.
- We can access elements using

$$p(i_{\# ext{extremities}} = 4, i_{\# ext{eyes}} = 2) = p(i_{\# ext{eyes}} = 2, i_{\# ext{extremities}} = 4)$$

Labelled Modes Tensors

Mode product

Let M, K, N be mutually disjoint mode sets. Given two tensors $x \in \mathbb{K}(M \cup K)$, $y \in \mathbb{K}(K \cup N)$, their mode product z := xy is the tensor $z \in \mathbb{K}(M \cup N)$ such that

$$z(i_M \times i_N) = \sum_{i_K} x(i_M \times i_K) y(i_K \times i_N).$$

Example

Recall example of tensor contraction. It now becomes

$$A \in \mathbb{K}(\{a, b, c\}), \quad B \in \mathbb{K}(\{b, d, e\}), \quad AB \in \mathbb{K}(\{a, c, d, e\})$$
$$(AB)(i_{\{a,c\}} \times i_{\{d,e\}}) = \sum_{i_b} A(i_{\{a,c\}} \times i_{\{b\}}) B(i_{\{b\}} \times i_{\{d,e\}}).$$

Tree Tensors

Mode tree

A triplet (V, E, D) where (V, E) is a tree and D a function mapping vertices to disjoint mode sets.

Tree tensor A tuple of tensors

$$x \in \bigotimes_{v \in V} \mathbb{K}(E(v) \cup D(v)).$$

Contracting the network is trivial:

$$x=\prod_{v\in V}x_v.$$

TreeTensors.jl

Create mode tree and tree tensor:

Contract network:

```
julia> prod(values(x.tensors)) # Lacks order
Tensor{Float64}([#= Modes :a,:b,:d,:c,:f,:e =#])
```

```
julia> contract(x) # Contract leaves-to-root
Tensor{Float64}([#= Modes :c,:a,:b,:f,:d,:e =#])
```

Conclusion

Labelled Modes (personal review)

- No notational blinkers.
- Shorter and more expressive formulae.
- Easy and bulletproof coding.

Try it yourself, it's free!

Conclusion

Material

- Slides: homepages.warwick.ac.uk/student/S.Etter/
- Tensors and TreeTensors packages: github.com/ettersi
- Publication:

S. Etter, Parallel ALS Algorithm for Solving Linear Systems in the Hierarchical Tucker Representation, SIAM Journal on Scientific Computing

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Mode Tags

Mode tag

A symbol t such that t(k) introduces a new mode based on k.

Row and Column mode tags

Tags R and C with special multiplication rules.

- R(k) only multiplies k and C(k) appearing on the left.
- C(k) only multiplies k and R(k) appearing on the right.

• If multiplication produces unpaired R(k)/C(k), rename to k. Write $D^2 := R(D) \cup C(D)$.

Example

- Operator $A \in \mathbb{K}(D^2)$, vector $x \in \mathbb{K}(D)$.
- Without last rule: $Ax \in \mathbb{K}(R(D))$.
- With last rule: $Ax \in \mathbb{K}(D)$.

Orthogonalisation

Maths paraphrased from [Ose11]

for
$$k = 1$$
 to $d - 1$ do
 $[q(\beta_k i_k; \beta_{k+1}), r(\beta_{k+1}, \alpha_{k+1})] := QR(x_k(\beta_k i_k; \alpha_{k+1}))$
 $x_k = q$
 $x_{k+1} := r \times_1 x_{k+1}$
end for

Maths in labelled modes notation



[Ose11] I. V. Oseledets. *Tensor Train Decomposition*. SIAM Journal on Scientific Computing (2011)

Orthogonalisation

Code

```
for (v,p) in edges(x, leaves_to_root)
    (q,r) = qr(x[v], PairSet(v,p))
    x[v] = q
    x[p] = r*x[p]
end
```

Truncation

Maths

for vertex v in root-to-leaves order do for child c of v do $(b, s_{v-c}, d) = SVD_{v-c}(x_v)$ $x_v = b \operatorname{diag}(s_{v-c})$ $x_c = \operatorname{diag}(s_{v-c}) d x_c$ Truncate s_{v-c} end for Truncate x_{ν} for child c of v do $x_{v} = \operatorname{diag}(s_{v-c}^{-1}) x_{v}$ end for end for

Truncation

Code

```
s = Dict{PairSet{Tree}, Tensor{real(scalartype(x))}}()
for (v,p) in edges_with_root(x, root_to_leaves)
    for c in children(v,p)
        e = PairSet(v,c)
        b,s[e],d = svd(x[v], e, maxrank())
        x[c] = scale(s[e],d)*x[c]
        x[v] = scale(b,s[e])
        s[e] = resize(s[e], Dict(e => rank(s[e])))
    end
    x[v] = resize(x[v])
        Dict(e => length(s[e]) for e in neighbor_edges(v)))
    x[v] = scale(x[v]).
        [1./s[e] for e in child_edges(v,p)]...)
end
```