# Labelled Modes <br> A New Notation for Tensor Networks 

Simon Etter

University of Warwick
GAMM Annual Meeting
7th March 2017

## Introduction

Mathematician's style goals

- Precise.
- Concise.
- General.

The tensor network notation dilemma

- Tensor trains: precise \& concise, but not general.
- Hierarchical Tucker: general, but not precise or concise.


## Introduction

## Furthermore

- Even hierarchical Tucker not really general: Requires rooted binary tree, free modes only in leaves.
- Even tensor trains not really precise and concise:

$$
x\left(i_{1}, \ldots, i_{d}\right)=\sum_{\alpha_{1}} \ldots \sum_{\alpha_{d-1}} x_{1}\left(i_{1}, \alpha_{1}\right) x_{2}\left(\alpha_{1}, i_{2}, \alpha_{2}\right) \ldots x_{d}\left(\alpha_{d-1}, i_{d}\right)
$$

- Coding is hard:

$$
\begin{aligned}
\mathrm{x}= & \mathrm{x}[1] \\
\mathrm{n}= & \mathrm{n}[1]
\end{aligned} \quad \begin{array}{rl}
\text { for } \mathrm{k}= & 2: \mathrm{d} \\
\mathrm{x}= & \mathrm{reshape}(\mathrm{x}, \quad(\mathrm{n}, \mathrm{r}[\mathrm{k}])) * \\
& \quad \operatorname{reshape}(\mathrm{x}[\mathrm{k}], \quad(\mathrm{r}[\mathrm{k}], \mathrm{n}[\mathrm{k}] * \mathrm{r}[\mathrm{k}+1])) \\
\mathrm{n} * & \mathrm{n}[\mathrm{k}]
\end{array}
$$

## Introduction

## Example

- Scientist counting fraction of species with particular features.
- Natural representation: $p(\#$ extremities $=x, \#$ eyes $=y, \ldots$ )
- Tensor representation: $p\left(i_{1}, \ldots, i_{d}\right)$


## Example

- Consider two tensors $A \in \mathbb{R}^{4 \times 5 \times 6}, B \in \mathbb{R}^{5 \times 7 \times 8}$.
- Only way to contract: 2 nd mode of $A$ with 1 st mode of $B$.
- Order of modes in result follows from mode sizes.
- Hence, $C=A B$ contains same information as

$$
C(a, c, d, e)=\sum_{b} A(a, b, c) B(b, d, e) .
$$

## Labelled Modes Tensors

## Generalised tuple

Let $D$ be a finite set and $\left(A_{k}\right)_{k \in D}$, a family of sets.

- A tuple $t \in X_{k \in D} A_{k}$ is a function

$$
t: D \rightarrow \bigcup_{k \in D} A_{k} \text { such that } t_{k} \in A_{k}
$$

- Traditional tuples: $D=\{1, \ldots, n\}$.


## Tensor

A tensor $x$ is a function mapping integer tuples to scalars,

$$
\mathbb{K}(D):=\left\{x: \underset{k \in D}{X}\left[n_{k}\right] \rightarrow \mathbb{K}\right\}, \quad \text { where } \quad\left[n_{k}\right]:=\left\{1, \ldots, n_{k}\right\}
$$

## Labelled Modes Tensors

## This construction is natural and intuitive!

Recall example of scientist counting species.

- Data naturally represented as $p \in \mathbb{R}(\{\#$ extremities, \#eyes $\})$.
- We can access elements using

$$
p\left(i_{\# \text { extremities }}=4, i_{\# \text { eyes }}=2\right)=p\left(i_{\# \text { eyes }}=2, i_{\# \text { extremities }}=4\right)
$$

## Labelled Modes Tensors

## Mode product

Let $M, K, N$ be mutually disjoint mode sets.
Given two tensors $x \in \mathbb{K}(M \cup K), y \in \mathbb{K}(K \cup N)$, their mode product $z:=x y$ is the tensor $z \in \mathbb{K}(M \cup N)$ such that

$$
z\left(i_{M} \times i_{N}\right)=\sum_{i_{K}} x\left(i_{M} \times i_{K}\right) y\left(i_{K} \times i_{N}\right)
$$

## Example

Recall example of tensor contraction. It now becomes

$$
\begin{gathered}
A \in \mathbb{K}(\{a, b, c\}), \quad B \in \mathbb{K}(\{b, d, e\}), \quad A B \in \mathbb{K}(\{a, c, d, e\}) \\
(A B)\left(i_{\{a, c\}} \times i_{\{d, e\}}\right)=\sum_{i_{b}} A\left(i_{\{a, c\}} \times i_{\{b\}}\right) B\left(i_{\{b\}} \times i_{\{d, e\}}\right)
\end{gathered}
$$

## Tree Tensors

Mode tree
A triplet $(V, E, D)$ where $(V, E)$ is a tree and $D$ a function mapping vertices to disjoint mode sets.

$$
\underset{\{c\}}{\{a\}} \underset{\{b\}}{\substack{\{c \mid}}\left\{\begin{array}{c}
\{d, g\}
\end{array}\right.
$$

## Tree tensor

A tuple of tensors

$$
x \in \underset{v \in V}{X} \mathbb{K}(E(v) \cup D(v))
$$

Contracting the network is trivial:

$$
x=\prod_{v \in V} x_{v}
$$

## TreeTensors.j

Create mode tree and tree tensor:

```
julia> using Tensors, TreeTensors;
julia> mtree = ModeTree([Mode(:c,2)],
    ModeTree([Mode(:a,2)]),
    ModeTree([Mode(:b,2)]),
    ModeTree([Mode(:d,2)],
        ModeTree([Mode(:e,2)]),
        ModeTree([Mode(:f,2)]),
    )
```

```
    );
```

    );
    julia> x = rand(mtree,2);

```
julia> x = rand(mtree,2);
```



Contract network:

```
julia> prod(values(x.tensors)) # Lacks order
Tensor{Float64}([#= Modes :a,:b,:d,:c,:f,:e =#])
julia> contract(x) # Contract leaves-to-root
Tensor{Float64}([#= Modes :c,:a,:b,:f,:d,:e =#])
```


## Conclusion

Labelled Modes (personal review)

- No notational blinkers.
- Shorter and more expressive formulae.
- Easy and bulletproof coding.

Try it yourself, it's free!

## Conclusion

## Material

- Slides: homepages.warwick.ac.uk/student/S.Etter/
- Tensors and TreeTensors packages: github.com/ettersi
- Publication:
S. Etter, Parallel ALS Algorithm for Solving Linear Systems in the Hierarchical Tucker Representation, SIAM Journal on Scientific Computing


## Acknowledgements

Thanks to Vladimir Kazeev, Robert Gantner, Christoph Schwab and Christoph Ortner for their help with this work!

## Mode Tags

## Mode tag

A symbol $t$ such that $t(k)$ introduces a new mode based on $k$.
Row and Column mode tags
Tags $R$ and $C$ with special multiplication rules.

- $R(k)$ only multiplies $k$ and $C(k)$ appearing on the left.
- $C(k)$ only multiplies $k$ and $R(k)$ appearing on the right.
- If multiplication produces unpaired $R(k) / C(k)$, rename to $k$.

Write $D^{2}:=R(D) \cup C(D)$.

## Example

- Operator $A \in \mathbb{K}\left(D^{2}\right)$, vector $x \in \mathbb{K}(D)$.
- Without last rule: $A x \in \mathbb{K}(R(D))$.
- With last rule: $\quad A x \in \mathbb{K}(D)$.


## Orthogonalisation

## Maths paraphrased from [Ose11]

for $k=1$ to $d-1$ do

$$
\begin{aligned}
& \qquad\left[q\left(\beta_{k} i_{k} ; \beta_{k+1}\right), r\left(\beta_{k+1}, \alpha_{k+1}\right)\right]:=\mathrm{QR}\left(x_{k}\left(\beta_{k} i_{k} ; \alpha_{k+1}\right)\right) \\
& \quad x_{k}=q \\
& \quad x_{k+1}:=r \times_{1} x_{k+1} \\
& \text { end for }
\end{aligned}
$$

Maths in labelled modes notation
for edge $v-p$ in leaves-to-root order do

$$
\begin{aligned}
& \quad(q, r):=\mathrm{QR}_{v-p}\left(x_{v}\right) \\
& \quad x_{v}:=q \\
& x_{p}:=r x_{p} \\
& \text { end for }
\end{aligned}
$$

[Ose11] I. V. Oseledets. Tensor Train Decomposition. SIAM Journal on Scientific Computing (2011)

## Orthogonalisation

## Code

```
for (v,p) in edges(x, leaves_to_root)
    (q,r) = qr (x[v], PairSet (v,p))
    x[v] = q
    x[p] = r*x[p]
end
```


## Truncation

## Maths

for vertex $v$ in root-to-leaves order do
for child $c$ of $v$ do

$$
\begin{aligned}
& \left(b, s_{v-c}, d\right)=\operatorname{SVD}_{v-c}\left(x_{v}\right) \\
& x_{v}=b \operatorname{diag}\left(s_{v-c}\right) \\
& x_{c}=\operatorname{diag}\left(s_{v-c}\right) d x_{c}
\end{aligned}
$$

$$
\text { Truncate } s_{v-c}
$$

end for
Truncate $x_{v}$
for child $c$ of $v$ do

$$
x_{v}=\operatorname{diag}\left(s_{v-c}^{-1}\right) x_{v}
$$

end for
end for

## Truncation

Code

```
s = Dict{PairSet{Tree}, Tensor{real(scalartype(x))}}()
for (v,p) in edges_with_root(x, root_to_leaves)
    for c in children(v,p)
            e = PairSet(v,c)
            b,s[e],d = svd(x[v], e, maxrank())
            x[c] = scale(s[e],d)*x[c]
            x[v] = scale(b,s[e])
            s[e] = resize(s[e], Dict(e => rank(s[e])))
    end
    x[v] = resize(x[v],
            Dict(e => length(s[e]) for e in neighbor_edges(v)))
    x[v] = scale(x[v],
            [1./s[e] for e in child_edges(v,p)]...)
end
```

